

# Digital Data Communication Techniques

## ITS323: Introduction to Data Communications

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ITS323Y11S1L06, Steve/Courses/ITS323/Lectures/digitaldata.tex, r1798

# Contents

Bit Errors

Error Detection

Error Correction

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# Bit Errors

- ▶ In digital transmission systems errors occur when a bit is altered between transmission and reception
- ▶ Single-bit errors
  - ▶ Only one bit altered, surrounding bits not affected
  - ▶ Caused by random noise
- ▶ Error burst
  - ▶ A group of bits near each other are affected (in error)
  - ▶ Caused by impulse noise or fading
  - ▶ Effects of burst errors are greater at higher data rates
- ▶ Require methods to detect errors, and correct where possible

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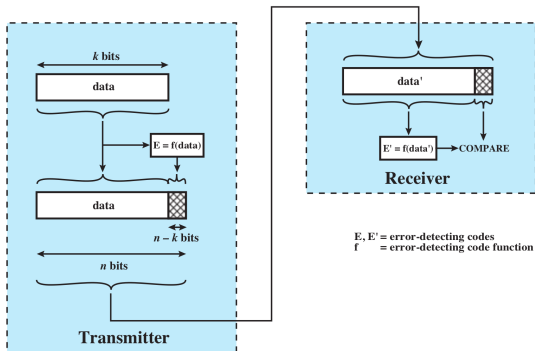
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# Error Detection

- ▶ Transmitter adds extra information to transmitted data, i.e. an **error-detecting code**
- ▶ Receiver recalculates the error-detecting code from received data, and compares to received error-detecting code
- ▶ If the same, good. If not, then error (in data or code). Still a chance that an error is not detected



## Error Detection with Parity Check

- ▶ Odd-parity check: append parity bit to block of data; resulting set of bits has odd number of ones
- ▶ Receiver detects an error if receiver bits has unexpected number of ones (transmitter and receiver both know parity scheme being used)

## Parity Check Example

Assume character  $S$  is to be sent using odd-parity check.  
What is transmitted? What happens if the last bit is corrupted? What about the last two bits?

# Error Detection with Cyclic Redundancy Check

- ▶ Parity checks are not good when multiple bit errors occur
- ▶ CRC is a powerful, commonly used error detection scheme
- ▶ Approach:
  - ▶  $k$  bits of data to send
  - ▶ Constant divisor known by transmitter/receiver,  $n - k + 1$  bits
  - ▶ Append  $n - k$  bits to data such that no remainder when divided by divisor
  - ▶ Transmit  $n$  bits
  - ▶ Receiver divides received  $n$  bits by divisor; if remainder, error detected
- ▶ Length and value of divisor is important for error detection capabilities (e.g. chance that one or more errors go undetected)
- ▶ CRC used in: Ethernet, HDLC, SATA, CDMA, PNG images, SD cards, ...



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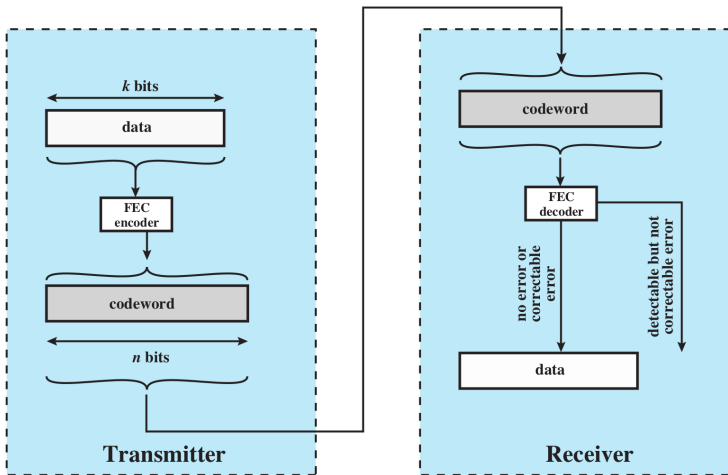
- ▶ What to do when error detected at receiver?
- ▶ Ask transmitter to send again, i.e. retransmit
  - ▶ Covered in Data Link Control Protocols lecture
  - ▶ Can be inadequate if link has high delay or many errors, e.g. wireless/satellite links
- ▶ **Forward Error Correction**: sender sends a codeword (instead of data); codeword chosen such that if error detected, receiver can **correct** the error without retransmission
- ▶ Depending on encoding scheme and pattern of errors, receiver may: detect and correct errors; detect, but not correct errors; not detect errors

# Error Correction

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# FEC with Hamming Distance

## Hamming Distance

- ▶ Number of bits of two  $n$ -bit sequences that differ
- ▶  $v_1 = 011011$ ,  $v_2 = 110001$ :  $d(v_1, v_2) = 3$

## Example FEC Encoder

- ▶ 2-bits of data mapped to 5-bit codeword ( $k = 2$ ,  $n = 5$ )

<i>Data</i>	<i>Codeword</i>
00	00000
01	00111
10	11001
11	11110

- ▶ If received codeword invalid, assume valid codeword that is unique minimum Hamming distance from received codeword was transmitted

# Error Correction Example 1

Data to send: 01; no transmission error

## Error Correction Example 2

Data to send: 01; 3rd bit transmitted is in error

## Error Correction Example 3

Data to send: 01; 1st and 4th bit transmitted in error

## Performance of Error Detection/Correction

- ▶ Aim to detect/correct as many errors as possible
- ▶ But error detection/correction require extra bits to be sent
- ▶  $k$  bits of useful data;  $n$  bits transmitted; efficiency  $\frac{k}{n}$
- ▶ Tradeoff: for a given amount of data,  $k$  bits
  - ▶ Increase  $n$ , more errors detected/corrected (GOOD)
  - ▶ Increase  $n$ , lower efficiency of transmission (BAD)