

Message Authentication Codes

CSS322: Security and Cryptography

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CSS322Y11S2L09, Steve/Courses/2011/S2/CSS322/Lectures/mac.tex, r2071

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MAC Algorithms

Attacks on Communications across Network

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2. Traffic analysis: encryption
3. Masquerade: message authentication
4. Content modification: message authentication
5. Sequence modification: message authentication
6. Timing modification: message authentication
7. Source repudiation: digital signatures
8. Destination repudiation: digital signatures

Message Authentication Functions

- ▶ Message authentication (and digital signature) mechanisms have two parts:
 1. Function that produces authenticator
 2. Protocol that enables receiver to verify authenticity
- ▶ Three types of authentication functions:
 1. Hash function
 2. Message encryption
 3. Message authentication code (MAC)

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Message Authentication Requirements and Functions

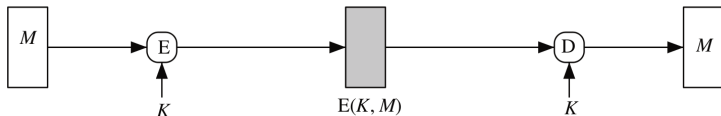
Authentication with Message Encryption

Authentication with Message Authentication Codes

Security of MACs

MAC Algorithms

Symmetric Encryption for Authentication



- ▶ Confidentiality: only B (and A) can recover plaintext
- ▶ Source Authentication: A is only other user with key; must have come from A
- ▶ Data Authentication: successfully decrypted; data has not been modified
- ▶ **Assumption**: decryptor can recognise correct plaintext

Recognising Correct Plaintext

Example 1

B receives ciphertext (supposedly from A , using shared secret key K):

DPNFCTEJLYONCJAEZRCLASJTDQFY

B decrypts with key K to obtain plaintext:

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- ▶ Was the plaintext encrypted with key K (and hence sent by A)?
- ▶ Is the ciphertext received the same as the ciphertext sent by A ?

Recognising Correct Plaintext

Example 2

B receives ciphertext (supposedly from A , using shared secret key K):

QEFPPQEBTOLKDJBPPXDBPLOOVX

B decrypts with key K to obtain plaintext:

FTUEUEFTQIDAZSYQEEMSQEADDKM

- ▶ Was the plaintext encrypted with key K (and hence sent by A)?
- ▶ Is the ciphertext received the same as the ciphertext sent by A ?

Recognising Correct Plaintext

Example 3

B receives ciphertext (supposedly from A , using shared secret key K):

0110100110101101010110111000010

B decrypts with key K to obtain plaintext:

0101110100001101001010100101110

- ▶ Was the plaintext encrypted with key K (and hence sent by A)?
- ▶ Is the ciphertext received the same as the ciphertext sent by A ?

Recognising Correct Plaintext

Example 1

- ▶ Assume the message is English
- ▶ Plaintext had expected structure; assume the plaintext is correct
 - ▶ Sent by A and has not been modified

Example 2

- ▶ Assume the message is English
- ▶ Plaintext had no structure in expected language; assume plaintext is incorrect
 - ▶ Either not sent by A or modified

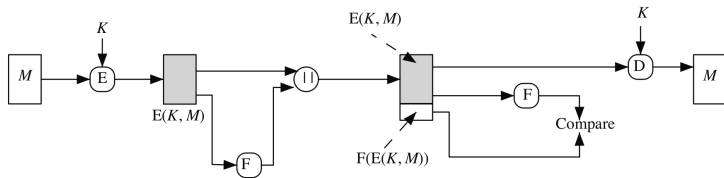
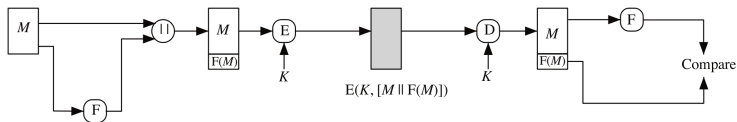
Example 3

- ▶ Binary data, e.g. image, compressed file
- ▶ Cannot know whether correct or incorrect

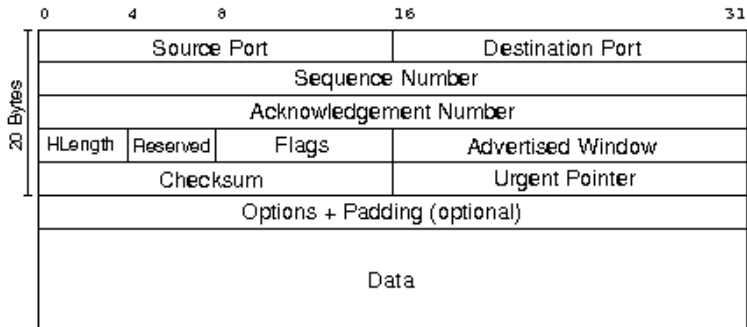
Recognising Correct Plaintext

- ▶ Valid plaintexts should be small subset of all possible messages
 - ▶ E.g. 26^n possible messages of length n ; only small subset are valid English phrases
- ▶ Plaintext messages have structure
- ▶ BUT automatically detecting structure can be difficult
- ▶ Add structure to make it easier, e.g.
 - ▶ Error detecting code or Frame Check Sequence
 - ▶ Packet header

Adding Error Detecting Code

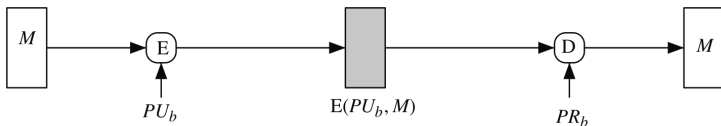


TCP Segment



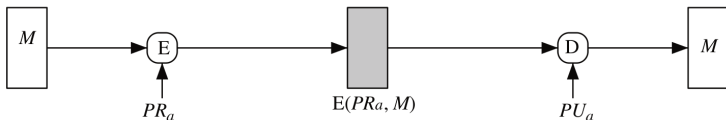
Public-Key Encryption for Authentication

- ▶ Only provides confidentiality
- ▶ Same assumption as before: plaintext must have structure so can be recognised as correct or incorrect



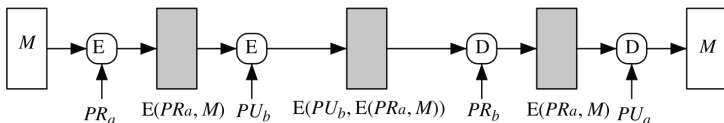
Public-Key Encryption for Authentication

- ▶ Data authentication (data has not been modified)
- ▶ **Digital signature**: proof that message came from A



Public-Key Encryption for Authentication

- ▶ Both confidentiality, authentication and digital signature



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- ▶ Append small, fixed-size block of data to message: cryptographic checksum or MAC

$$T = \text{MAC}(K, M)$$

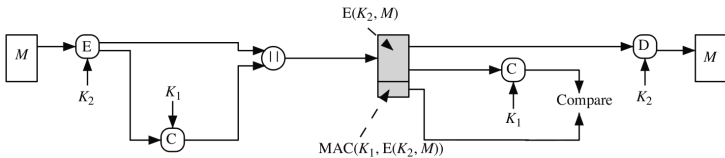
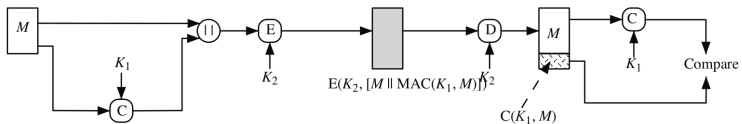
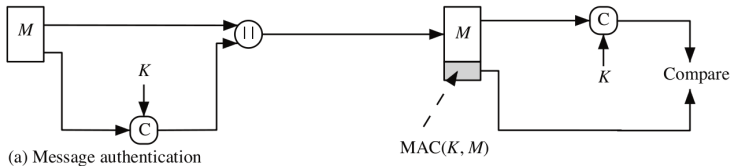
M = input message

MAC = MAC function

K = shared secret key of k bits

T = message authentication code (or tag) of n bits

- ▶ MAC function also called *keyed hash function*
- ▶ MAC function similar to encryption, but does not need to be reversible
 - ▶ Easier to design stronger MAC functions than encryption functions



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Requirement of MACs

Objective of Attacker

- ▶ Assume MAC function is known, key K is not
- ▶ For valid MAC code for given message x

Requirement of MAC Function

Computation Resistance : given one or more text-MAC pairs $[x_i, MAC(K, x_i)]$, computationally infeasible to computer any text-MAC pair $[x, MAC(K, x)]$ for new input $x \neq x_i$

Security of MACs

Brute Force Attack on Key

- ▶ Attacker knows $[x_1, T_1]$ where $T_1 = \text{MAC}(K, x_1)$
- ▶ Key size of k bits: brute force on key, 2^k
- ▶ But ... many tags match T_1
- ▶ For keys that produce tag T_1 , try again with $[x_2, T_2]$
- ▶ Effort to find K is approximately 2^k

Brute Force Attack on MAC value

- ▶ For x_m , find T_m without knowing K
- ▶ Similar effort required as one-way/weak collision resistant property for hash functions
- ▶ For n bit MAC value length, effort is 2^n

Effort to break MAC: $\min(2^k, 2^n)$

Security of MACs

Cryptanalysis

- ▶ Many different MAC algorithms; attacks specific to algorithms
- ▶ MAC algorithms generally considered secure

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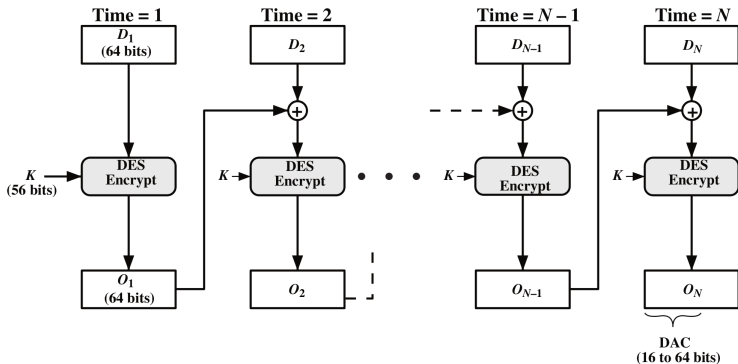
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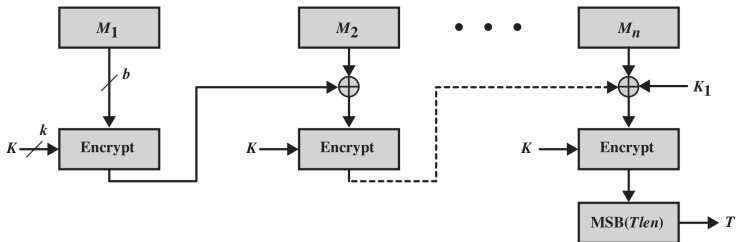
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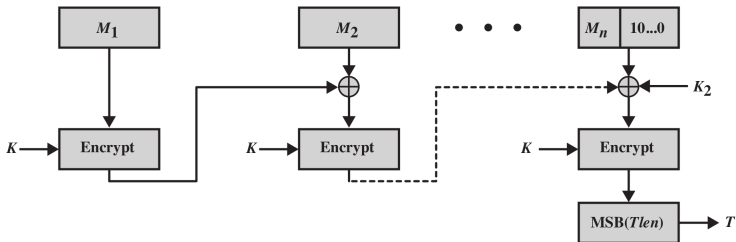
MACs Based on Block Ciphers

- ▶ Data Authentication Algorithm (DAA): based on DES; considered insecure
- ▶ Cipher-Based Message Authentication Code (CMAC): mode of operation used with Triple-DES and AES
- ▶ OMAC, PMAC, UMAC, VMAC, ...





(a) Message length is integer multiple of block size



(b) Message length is not integer multiple of block size

HMAC

- ▶ MAC function derived from cryptographic hash functions
 - ▶ MD5/SHA are fast in software (compared to block ciphers)
 - ▶ Libraries for hash functions widely available

$$\text{HMAC}(K, M) = H((K \oplus \text{opad}) || H((K \oplus \text{ipad}) || M))$$

where $\text{ipad} = 00110110$ repeated, $\text{opad} = 01011100$ repeated

- ▶ Security of HMAC depends on security of hash function used